<u>**Amendments to the Claims:**</u>

This listing of claims will replace all prior versions, and listings, of claims in the application:

<u>**Listing of Claims:**</u>

Claims 1-42.   (Canceled).

43.   (Currently Amended) A method of rate adaptation in a communication apparatus, comprising:

~~providing information bits of a prescribed data rate to an input of an encoder, the encoder having a prescribed code rate;~~

~~adapting the prescribed code rate of the encoder and providing coded bits, based on an adapted code rate, to an output of the encoder, the prescribed code rate being adapted to the adapted code rate for providing a coding gain; and~~

<u>encoding information bits by an encoder at a coding rate, the coding rate being fixed to a specific value among at least two different values using a puncturing block in the encoder according to a ratio of a channel interleaver size and a number of the information bits;</u>

performing repetition or puncturing of the coded bits ~~by a rate matching device,~~
~~the output of the encoder being coupled to an input of the rate matching device~~ for matching a
size of the coded bits to the channel interleaver size; and

interleaving the result of the repetition or puncturing ~~which is provided from an~~
~~output of the rate matching device to an input of an interleaver~~ according to the channel
interleaver size.

44. (Canceled).

45. (Currently Amended) The method of claim 43, wherein ~~the adapted code rate of~~
~~the encoder is one of~~ the at leaset two different values included at least two of 1/2, 1/3, 1/4,
and 1/5.

46. (Canceled).

47. (Currently Amended) The method of claim 43, wherein the encoder is a turbo
encoder ~~with a maximum code rate of 1/5~~.

48. (Canceled).

49.   (Previously Presented)   The method of claim 43, wherein symbol puncturing is enabled for symbol groups having indices 2j and 2j+1 if (j●k) mod J < K, wherein 'I' is a number of information bits per frame, 'J' equals $\lfloor I/2 \rfloor$, 'N' is a size of the interleaver, 'K' equals $\lfloor (L-N)/2 \rfloor$, and 'L' is a number of coded bits, and wherein each of the symbol groups comprises L/I coded bits.

50.   (Previously Presented)   The method of claim 49, wherein the information bits include data bits and a pattern used to puncture the symbol group 'i' for the adapted code rate of 1/3 turbo code rate when a prescribed ratio < N ≤ 3I is given by $P_{(i \bmod 2)}$, wherein 'i' is an index of the symbol groups and ranges from 0 to I-1, and wherein the pattern to puncture symbol groups corresponding to coded bits of data bits is '110' for P_0 and '101' for P_1, where '1' indicates no puncturing of the coded bit in the symbol group 'i' and '0' indicates puncturing of the coded bit in the symbol group 'i'.

51.   (Previously Presented)   The method of claim 50, wherein the information bits further includes tail bits and a pattern to puncture symbol groups corresponding to coded bits of tail bits is '101' for P_0 and '101' for P_1.

52. (Previously Presented) The method of claim 49, wherein the information bits include data bits and a pattern used to puncture the symbol group 'i' for the adapted code rate of 1/4 turbo code rate when $3I < N \leq 4I$ is given by $P_{(i \bmod 2)}$, wherein 'i' is an index of the symbol groups and ranges from 0 to I-1, and wherein the pattern to puncture symbol groups corresponding to coded bits of data bits is '1011' for P_0 and '1110' for P_1, where '1' indicates no puncturing of the coded bit in the symbol group 'i' and '0' indicates puncturing of the coded bit in the symbol group 'i'.

53. (Previously Presented) The method of claim 52, wherein the information bits further include tail bits and a pattern to puncture symbol groups corresponding to coded bits of tail bits is '1011' for P_0 and '1011' for P_1.

54. (Previously Presented) The method of claim 49, wherein the information bits include data bits and a pattern used to puncture the symbol group 'i' for the adapted code rate of 1/5 turbo code rate when $4I < N \ [[\leq]] < 5I$ is given by $P_{(i \bmod 2)}$, wherein 'i' is an index of the symbol groups and ranges from 0 to I-1, and wherein the pattern to puncture symbol groups corresponding to coded bits of data bits is '11101' for P_0 and '11011' for P_1, where '1' indicates no puncturing of the coded bit in the symbol group 'i' and '0' indicates puncturing of the coded bit in the symbol group 'i'.

55.   (Previously Presented)   The method of claim 54, wherein the information bits further include tail bits and a pattern to puncture symbol groups corresponding to coded bits of tail bits is '11011' for P_0 and '11011' for P_1.

56 -83. (Canceled).

84.   (Previously Presented)   The method of claim 43, wherein the method is implemented during variable data rate mode and/or flexible data rate mode.

85.   (Previously Presented)   The method of claim 43, wherein the method is used for radio configuration (RC)4 of a physical channel.

86 - 103. (Canceled).